

AN ASYMPTOTIC ESTIMATE OF THE NUMBER OF SOLUTIONS
OF A SPECIAL SYSTEM OF BOOLEAN EQUATIONS

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In this paper a special class of systems of Boolean equations is investigated. For a “typical” case of such systems an asymptotic estimate for the number of solutions is determined.

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Introduction. Many problems of discrete mathematics, including the problems which are traditionally considered to be complex, lead to the solutions of the systems of Boolean equations of the form

$$\begin{cases} f_i(x_1, \dots, x_n) = 0, \\ i = 1, \dots, l, \end{cases} \quad (1)$$

or to the revealing of those conditions, under which the system (1) has a solution. In general, the problem of realizing whether the system (1) has a solution or not is NP-complete [1]. Therefore, it is often necessary to explore a number of solutions in the “typical” case, or to consider special classes of systems of equations, using their specificity.

Some Necessary Definitions. Let $\{M(n)\}_{n=1}^{\infty}$ be the collection of sets such that $|M(n)| \xrightarrow{n \rightarrow \infty} \infty$ ($|M|$ is the power of the set M), and $M^S(n)$ be the subset of all elements from $M(n)$, which have a property S . We say that almost all elements of the set $M(n)$ have a property S , if $|M^S(n)| / |M(n)| \xrightarrow{n \rightarrow \infty} 1$.

Everywhere below the notation \log stands for the logarithm of the base 2; $[a]$ denotes the integer part of a .

We denote by $S_{n,l}$ the set of all the systems of the form (1), where $f_i(x_1, \dots, x_n)$, $i = 1, \dots, l$, are pairwise different Boolean functions in variables x_1, x_2, \dots, x_n . It is easy to see that $|S_{n,l}| = C_{2^n}^l$.

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Let $B = \{0, 1\}$, $B^n = \{\tilde{\alpha}/\tilde{\alpha} = (\alpha_1, \alpha_2, \dots, \alpha_n), \alpha_i \in B, 1 \leq i \leq n\}$. The vector $\tilde{\alpha}_i = (\alpha_1, \alpha_2, \dots, \alpha_n) \in B^n$ is called a solution of (1), if

$$\begin{cases} f_i(\alpha_1, \alpha_2, \dots, \alpha_n) = 0, \\ i = 1, \dots, l. \end{cases}$$

We denote by $t(S)$ the number of solutions of the system S . In [2, 3] it is found the asymptotics of the number of solutions $t(S)$ for almost all the systems S of the set $S_{n,l}$ for whole range values of parameter l , as $n \rightarrow \infty$. Moreover, in [3] a general case, the systems of equations in k -valued logic was considered. In this work a class of systems of equations of the special form is considered. The asymptotic behavior of the number of solutions of equations systems in a “typical” case is found.

1. Let $L(n)$ denote the set of all linear Boolean functions, depending on the variables x_1, x_2, \dots, x_n , i.e. $L(n) = \{c_0 + c_1 x_1 + \dots + c_n x_n / c_i = 0, 1; i = 1, \dots, n\}$, where $+$ indicates addition modulo.

2. Obviously, $|L(n)| = 2^{n+1}$. Let $f_i = f_i(x_1, \dots, x_n) = L_{i1} \dots L_{im}$, $i = 1, \dots, l$, where $L_{ij} = L_{ij}(x_1, \dots, x_n) \in L(n)$, $j = 1, \dots, m$.

Then the system of equations (1) takes the form

$$\begin{cases} L_{i1} \dots L_{im} = 0, \\ i = 1, \dots, l. \end{cases} \quad (2)$$

Let $R(n, l, m)$ denote the set of all systems of l equations of the form (2) ($f_i \neq f_j$, if $i \neq j$) and let $c(n, m) = C_{2^{n+1}}^m$. It is easy to see that $|R(n, l, m)| = C_{c(n, m)}^l$. For the numbers of the solutions $t(S)$ of almost all systems S of $R(n, l, m)$ the following statement is true (here and further $f(n) \sim g(n)$ stands for the relation $f(n)/g(n) \xrightarrow{n \rightarrow \infty} 1$, $f(n) = o(g(n))$ for $f(n)/g(n) \xrightarrow{n \rightarrow \infty} 0$).

Theorem.

1. If $n - \ell(m - \log(2^m - 1)) \xrightarrow{n \rightarrow \infty} \infty$ and $l^2 = o(c(n, m))$, $m^2 = o(2^n)$, then for almost all systems S of $R(n, l, m)$ we have $t(S) \sim 2^n(1 - 2^{-m})^l$.

2. If $n - \ell(m - \log(2^m - 1)) \xrightarrow{n \rightarrow \infty} -\infty$ and $m^2 = o(2^n)$, then almost all systems S of $R(n, l, m)$ have no solutions.

3. If $n - \ell(m - \log(2^m - 1))$ is bounded as $n \rightarrow \infty$ and $l^2 = o(c(n, m))$, $m^2 = o(2^n)$, then for almost all systems of $R(n, l, m)$ the number of solutions $t(S)$ is bounded from above by an arbitrary function $\varphi(n)$, satisfying the condition $\varphi(n) \xrightarrow{n \rightarrow \infty} \infty$.

Proof. The following known or easily checking inequalities hold:

1. The first Chebyshev inequality [4]. Let the random variable ξ take the non-negative values and have mathematical expectation $M\xi$. Then for any $t > 0$ we have

$$P(\xi \geq t) \leq M\xi/t.$$

2. The second Chebyshev inequality [4]. Let the above random variable ξ have a dispersion $D\xi$. Then for any $t > 0$ we have

$$P(|\xi - M\xi| \geq t) \leq D\xi/t^2.$$

3. For any $x > 1$

$$(1 - 1/x)^x < e^{-1}.$$

4. For any $x \geq 1$

$$(1 - 1/x)^x > e^{-1 - \frac{1}{x}}.$$

5. For any natural n with $1 \leq m \leq n$ we have

$$C_n^m < \left(\frac{en}{m}\right)^m.$$

6. Let $b(k; n, p) = C_n^k p^k q^{n-k}$, where $0 < p, q < 1, p + q = 1$. Then for $r > np$

$$\sum_{j=0}^{n-r} b(r+j; n, p) < b(r; n, p)(r+1)q/(r+1 - (n+1)p)$$

(the estimate of the “tail” of the binomial distribution [4]).

7. Bernoulli inequality: for any $x \geq -1$ and natural n , we have

$$(1+x)^n \geq 1+nx.$$

Consider another set $R'(n, l, m)$ of the systems of l equations of the form (2), where $L_{ij} = L_{ij}(x_1, \dots, x_n) \in L(n), i = 1, \dots, l, j = 1, \dots, m$ (a condition $f_i \neq f_j$ when $i \neq j$ is canceled). In addition, the systems of equations are assumed to be “ordered”, i.e. two systems S' and S'' from $R'(n, l, m)$ are different, if one of them is obtained from other by a transposition of non-equivalent equations. It's easy to observe that $|R'(n, l, m)| = c(n, m)^l$.

Let S be a system in $R(n, l, m)$. Arranging (transpositions by all the variations) the equations in S , we obtain $l!$ new systems, differing from each other by the transposition of the equations. Thus, from the set $R(n, l, m)$ we obtain a new set $R''(n, l, m)$ of ordered and non-repetitive (containing no pair of equivalent equations) systems. It is evident that $|R''(n, l, m)| = |R(n, l, m)|l!$. Using Bernoulli inequality, one can estimate the relation

$$\begin{aligned} \frac{|R''(n, l, m)|}{|R'(n, l, m)|} &= \frac{l! \cdot |R(n, l, m)|}{(c(n, m))^l} = \\ &= \frac{l! \cdot C_{c(n, m)}^l}{(c(n, m))^l} = \frac{c(n, m) \cdot (c(n, m) - 1) \times \dots \times (c(n, m) - l + 1)}{(c(n, m))^l} = \\ &= \prod_{i=0}^{l-1} \left(1 - \frac{i}{c(n, m)}\right) \geq \left(1 - \frac{l-1}{c(n, m)}\right)^l \geq 1 - \frac{(l-1)}{c(n, m)}l \rightarrow 1, \end{aligned}$$

if $l^2/c(n, m) \xrightarrow{n \rightarrow \infty} \infty$. Thus, if $l^2 = o(c(n, m))$, then any statement concerning the almost all systems of the set $R'(n, l, m)$ is also true for almost all the systems of the set $R''(n, l, m)$.

Let almost all systems of $R''(n, l, m)$ have property E , which is invariant with respect to the transpositions of the equations of system. It is easy to see that almost all the systems of the set $R(n, l, m)$ will also have property E . Thus, for the Proof of the Theorem when $l^2 = o(c(n, m))$ it will be enough to consider the set $R'(n, l, m)$ instead of $R(n, l, m)$.

Next, let S be a system of $R'(n, l, m)$. Arranging (transpositions by all the variations) linear functions in the equations of the system, we obtain $(m!)^l$ new

ordered systems, different from each other by permutation of linear functions in the equations. Thus, from the set of $R'(n, l, m)$ a new set of $R'''(n, l, m)$ ordered systems can be obtained. It is clear that $|R'''(n, l, m)(n, l, m)| = (m!)^l \cdot |R'(n, l, m)|$, i.e.

$$|R'''(n, l, m)(n, k, p)| = (2^{n+1} \cdot (2^{n+1} - 1) \times \dots \times (2^{n+1} - m + 1))^l. \quad (3)$$

Finally, we denote by $R''''(n, l, m)$ the expansion of the set $R'''(n, l, m)$ (different systems of $R''''(n, l, m)$ are allowed to have same equations). It is easy to see that

$$|R''''(n, l, m)| = (2^{n+1})^{lm}. \quad (4)$$

From (3), (4), using Bernoulli inequality, we obtain

$$\begin{aligned} \frac{|R'''(n, l, m)|}{|R''''(n, l, m)|} &= \frac{(2^{n+1} \cdot (2^{n+1} - 1) \times \dots \times (2^{n+1} - m + 1))^l}{(2^{n+1})^{lm}} = \\ &= \prod_{i=0}^{m-1} \left(1 - \frac{i}{2^{n+1}}\right) \geq \prod_{i=0}^{m-1} \left(1 - \frac{m-1}{2^{n+1}}\right) \geq 1 - \frac{m(m-1)}{2^{n+1}} \rightarrow 1, \end{aligned}$$

as $m^2 = o(2^n)$ ($n \rightarrow \infty$). Thus, if $m^2 = o(2^n)$, then any assertion concerning the almost all systems of set $R''''(n, l, m)$ is true for almost all systems of $R'''(n, l, m)$.

We consider $R''''(n, l, m)$ as a space of events, where every event $S \in R''''(n, l, m)$ takes place with the probability $1/R''''(n, l, m) = 2^{-(n+1)ml}$. Consider the random value $\xi_S(\tilde{\alpha})$, which is connected with $S \in R''''(n, l, m)$ as follows: $\xi_S(\tilde{\alpha}) = 1$, if $\tilde{\alpha}$ is the solution of the system S, otherwise $\xi_S(\tilde{\alpha}) = 0$.

From the definition and properties of a linear function it follows that the number of equations of the form $L_i^1 \& L_i^2 \& \dots \& L_i^m = 0$, for which $\tilde{\alpha}$ is a solution, is equal to $(2^{(n+1)m} - 2^{nm})^l = 2^{nml} (2^m - 1)^l$.

From this and (4) it follows that $P(\xi_S(\tilde{\alpha}) = 1) = (1 - 1/2^m)^l$.

Consider another random value $v = \sum_{\tilde{\alpha} \in B^n} \xi_S(\tilde{\alpha})$. Random value v has a binomial distribution, because

$$p(v = j) = C_{2^n}^j \left(1 - \frac{1}{2^m}\right)^{lj} \left(1 - \left(1 - \frac{1}{2^m}\right)^l\right)^{2^n - j}.$$

Hence, $Mv = 2^n (1 - 1/2^m)^l$ and $Dv = 2^n (1 - 1/2^m)^l \cdot (1 - (1 - 1/2^m)^l)$, where Mv and Dv are the random value v mathematical expectation and dispersion respectively.

Let $n - \ell(m - \log(2^m - 1)) \xrightarrow{n \rightarrow \infty} \infty$. It means that $Mv = 2^n (1 - 1/2^m)^l \xrightarrow{n \rightarrow \infty} \infty$. Using the Chebishev's second inequation when $t = Mv/g(n)$, where $g(n)$ is an arbitrary function, satisfying the conditions $g(n) \xrightarrow{n \rightarrow \infty} \infty$, $g^2(n) = o(Mv)$, we obtain $P\left(|v - Mv| > \frac{Mv}{g(n)}\right) < \frac{g^2(n)}{Mv} \rightarrow 0$ as $n \rightarrow \infty$. Hence and from the definition of random value v it follows that almost all systems of the set $R''''(n, l, m)$ have a number of solutions, which asymptotically equals Mv . Therefore, under the conditions of the theorem, almost all systems of $R(n, l, m)$ have also number of solutions asymptotically equal $Mv = 2^n (1 - 1/2^m)^l$.

The first statement of the Theorem is proved.

Let $n - \ell(m - \log(2^m - 1)) \xrightarrow{n \rightarrow \infty} -\infty$. Then $Mv = 2^n (1 - 1/2^m)^l \xrightarrow{n \rightarrow \infty} 0$. Using Chebishev's first inequation when $t = l$, we obtain $P(v \geq 1) \xrightarrow{n \rightarrow \infty} 0$ and, therefore, $P(v = 0) \xrightarrow{n \rightarrow \infty} 1$. Hence, it follows that almost all systems S of the set $R'''(n, l, m)$ have no solution. Therefore, when $m^2 = o(2^n)$, $l^2 = o(c(n, m))$ the second statement of the Theorem is proved.

It is easy to see that it is enough to require $m^2 = o(2^n)$, because for the greater values of the parameter l the statement of the Theorem also is true (the number of solutions of the system does not increase together with the number of equations).

Now let $n - \ell(m - \log(2^m - 1))$ is bounded, as $n \rightarrow \infty$. Then $Mv = 2^n (1 - 1/2^m)^l$ is also bounded as $n \rightarrow \infty$. Using the inequations 6 and 5, we obtain

$$\begin{aligned} P(v > r) &= \sum_{i=0}^{2^n-r} c(n, r+i) \left(1 - \frac{1}{2^m}\right)^{l(r+i)} \left(1 - \left(1 - \frac{1}{2^m}\right)^l\right)^{2^n-r-i} < \\ &< c(n, r) \left(1 - \frac{1}{2^m}\right)^{lr} \left(1 - \left(1 - \frac{1}{2^m}\right)^l\right)^{2^n-r} (r+1) \times \\ &\times \left(1 - \left(1 - \frac{1}{2^m}\right)^l\right) / \left(r+1 - (2^n+1) \left(1 - \frac{1}{2^m}\right)^l\right) < \left(\frac{eMv}{r}\right) \rightarrow 0, \end{aligned}$$

as $r \rightarrow \infty$, because Mv is bounded. Therefore, for almost all systems of the set $R'''(n, l, m)$ and, so, for almost all systems of the set $R(n, l, m)$ the third statement of the Theorem holds. \square

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